

A new generalization of Mantel's theorem to k -graphs

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Abstract

Let the k -graph Fan^k consist of k edges that pairwise intersect exactly in one vertex x , plus one more edge intersecting each of these edges in a vertex different from x . We prove that, for n sufficiently large, the maximum number of edges in an n -vertex k -graph containing no copy of Fan^k is $\prod_{i=1}^k \lfloor \frac{n+i-1}{k} \rfloor$, which equals the number of edges in a complete k -partite k -graph with almost equal parts. This is the only extremal example. This result is a special case of our more general theorem that applies to a larger class of excluded configurations.

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1. Introduction

The first theorem in extremal graph theory is Mantel's 1907 result, which determines the maximum number of edges in a triangle-free graph on n vertices (cf. Turán [22]). There are several possible generalizations of this problem to k -uniform hypergraphs (k -graphs for short). One was suggested by Katona [9] and Bollobás [1] (see Frankl–Füredi [4,5], de Caen [2], Sidorenko [20], Shearer [19], Keevash–Mubayi [10], Pikhurko [16]). Another extension, the so-called *expanded triangle*, was studied by Frankl [3] and Keevash–Sudakov [11]. In this paper we provide yet another generalization.

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Let Fan^k be the k -graph comprising $k + 1$ edges E_1, \dots, E_k, E , with $E_i \cap E_j = \{x\}$ for all $i \neq j$, where $x \notin E$, and $|E_i \cap E| = 1$ for all i . In other words, k edges share a single common vertex x and the last edge intersects each of the other edges in a single vertex different from x . Please note that Fan^2 is simply a triangle, and in this sense Fan^k generalizes the definition of K_3 . There is another, perhaps more subtle way that Fan^k is an extension of K_3 .

Call a hypergraph *simple* if every two edges share at most one vertex. One of the formulations of the celebrated Erdős–Faber–Lovász conjecture states that the minimum number of edges in a simple k -graph that is not k -partite is $k + 1$. Kahn [8] proved this with $k + 1$ replaced by $(1 + o(1))k$, but the question of the exact value remains open. If the conjecture is true, then Fan^k is a simple k -graph that is not k -partite with the minimum number of edges, and in this sense it generalizes a 2-graph triangle.

For $l \geq k$, let $T_l^k(n)$ be the complete l -partite k -graph with part sizes $\lfloor n/l \rfloor$ or $\lceil n/l \rceil$: every edge of $T_l^k(n)$ has at most one vertex in each of the l parts, and all edges subject to this restriction are present. Let

$$t_l^k(n) = |T_l^k(n)|.$$

(We identify a k -graph with its edge set.) It is convenient to agree that $T_l^k(n) = \emptyset$ and $t_l^k(n) = 0$ if $l < k$. Given a k -graph F , we write $\text{ex}(n, F)$ for the maximum number of edges in an n -vertex k -graph containing no copy of F . Mantel proved that $\text{ex}(n, \text{Fan}^2) = t_2^2(n)$ for all positive n . Here we generalize this to $k > 2$, for large n .

Theorem 1. *Let $k \geq 3$. Then, for all sufficiently large n , the maximum number of edges in an n -vertex k -graph containing no copy of Fan^k is $t_k^k(n) = \prod_{i=1}^k \lfloor \frac{n+i-1}{k} \rfloor$. The only k -graph for which equality holds is $T_k^k(n)$.*

Our approach to proving Theorem 1 comes from two recent papers by the current authors [13,15]. Although the paper [15] has been accepted by the Journal of Combinatorial Theory, Series B, its publication is suspended for an indefinite period of time because of a disagreement over the copyright between the author and the publisher. We feel that the approach is quite versatile and may be applicable to other hypergraph Turán problems. Therefore, we give a complete description of the method and provide self-contained proofs for any claims from [15].

So suppose that we wish to prove that $\text{ex}(n, F) = t_l^k(n)$ for a given F . The method has four steps:

Step 1. Define an appropriately chosen family \mathcal{F} of k -graphs such that $F \in \mathcal{F}$. There is no general recipe for \mathcal{F} . A particular property that \mathcal{F} should possess is that any F -free k -graph of order n can be made \mathcal{F} -free by removing $o(n^k)$ edges. Then $\text{ex}(n, F) = \text{ex}(n, \mathcal{F}) + o(n^k)$ but, hopefully, $\text{ex}(n, \mathcal{F})$ is easier to analyze.

Step 2. Prove that \mathcal{F} is *stable* with respect to $T_l^k(n)$. Loosely speaking, this means that every \mathcal{F} -free k -graph G on n vertices with close to $\text{ex}(n, \mathcal{F})$ edges can be transformed to $T_l^k(n)$ without changing too many edges.

Step 3. From the stability of \mathcal{F} , deduce the stability of F . (We use the property of \mathcal{F} from Step 1, whose proof is combined with Step 3 in this article.)

Step 4. Using the stability of F , deduce the exact result $\text{ex}(n, F) = t_l^k(n)$. This technique was first employed by Simonovits [21] to determine $\text{ex}(n, F)$ exactly for color-critical 2-graphs F . Recently, stability has been used to determine exact results for several hypergraph Turán problems [6,10–12,15,16].

The next three sections give the details of Steps 2–4, culminating in a proof of Theorem 1. Actually, our main result, Theorem 3 proved in Section 4, determines the exact extremal function for a more general configuration which includes Fan^k as a special case. We next define the family used in Step 1.

Fix $l \geq k \geq 2$. Let \mathcal{F}_l^k be the set of all minimal k -graphs F such that there is an l -set C , called the *core*, such that at least one edge $D \in F$ lies entirely in C and every pair of vertices of C is covered by an edge of F . (Of course, it suffices to consider only pairs not inside D .) Let F_l^k be the k -graph with edges: $[k]$ and $E_{ij} \cup \{i, j\}$ over all pairs $\{i, j\} \in \binom{[l]}{2} \setminus \binom{[k]}{2}$, where E_{ij} are pairwise disjoint $(k - 2)$ -sets disjoint from $[l]$. Clearly, $F_l^k \in \mathcal{F}_l^k$. Note that

- $\mathcal{F}_k^k = \{F_k^k\}$ and F_k^k is the k -graph of one edge,
- $\mathcal{F}_l^2 = \{K_l^2\}$,
- $F_{k+1}^k = \text{Fan}^k$.

For $l \geq k \geq 2$ a particular \mathcal{F}_{l+1}^k -free k -graph is $T_l^k(n)$. It is easy to see this, since if $T_l^k(n)$ contains a copy of $F \in \mathcal{F}_{l+1}^k$, then the vertex set in $T_l^k(n)$ playing the role of C must have at most one point in each part of $T_l^k(n)$ but there are not enough parts to accommodate these $l + 1$ vertices. Consequently, the maximum size of an n -vertex \mathcal{F}_{l+1}^k -free k -graph is at least $t_l^k(n)$. In fact, we have an equality:

Theorem 2. Let $n \geq l \geq k \geq 2$, and let G be an n -vertex \mathcal{F}_{l+1}^k -free k -graph. Then $|G| \leq t_l^k(n)$, and if equality holds then $G = T_l^k(n)$.

This result can be proved by a straightforward modification of the proof of Theorem 1 in [13]. Also, one can obtain it as a by-product of our proof of Theorem 4 below (see the remark following the inequality (4)).

The main theorem of the current paper is the following extension of Theorem 1.

Theorem 3. Let $l \geq k \geq 2$. Then, for all sufficiently large n , we have $\text{ex}(n, F_{l+1}^k) = t_l^k(n)$ and $T_l^k(n)$ is the unique maximum F_{l+1}^k -free k -graph of order n .

Let us specify here the notation we are going to use. We write $V(G)$ for the vertex set of a k -graph G . Given a vertex $x \in V(G)$, the *link* of x is the $(k - 1)$ -graph

$$L_G(x) = \{S \setminus \{x\} : S \in G, S \ni x\},$$

and the *degree* is $\text{deg}_G(x) = |L_G(x)|$. The *codegree* of x and y , written $\text{codeg}_G(x, y)$, is the number of edges in G containing both x and y , and the *neighborhood* of x is

$$N_G(x) = \{y : \text{codeg}(x, y) > 0, y \neq x\}.$$

Given $X \subset V(G)$, let $e_G(X)$ be the number of edges in G that contain at least two vertices from X . In all cases above, we omit the subscript G if the k -graph G is obvious from context.

For $S \subset V(G)$, we write $G[S]$ for the hypergraph induced by G on S . Two k -graphs F and G of the same order are m -close if we can add or remove at most m edges from the first graph and make it isomorphic to the second; in other words, for some bijection $\sigma : V(F) \rightarrow V(G)$ the symmetric difference between $\sigma(F) = \{\sigma(D) : D \in F\}$ and G has at most m edges.

The notation $a \pm b$ means a number between $a - b$ and $a + b$.

2. Step 2: \mathcal{F}_{l+1}^k is stable

Our goal in this section is to prove the following stability result.

Theorem 4. *For any $l \geq k \geq 2$ and $\delta > 0$ there exist $\varepsilon > 0$ and M such that the following holds for all $n > M$: If G is an n -vertex \mathcal{F}_{l+1}^k -free k -graph with at least $t_l^k(n) - \varepsilon n^k$ edges, then G is δn^k -close to $T_l^k(n)$.*

The proof of Theorem 4 has many similarities to that in [13, Theorem 3]. Thus we will refer to [13] for proofs of some claims, when the arguments are identical. In particular, we use the following facts shown in [13].

Equation (1) in [13]. For any $l \geq k \geq 2$ and $0 \leq s \leq n$ we have

$$t_{l-1}^k(n-s) + s \cdot t_{l-1}^{k-1}(n-s) \leq t_l^k(n). \tag{1}$$

Hint. The left-hand side of (1) is the number of edges in the complete l -partite k -graph with one part of size s and other part sizes being $\lfloor \frac{n-s}{l-1} \rfloor$ and $\lceil \frac{n-s}{l-1} \rceil$.

Claim 1 in [13]. For any $l \geq k \geq 2$ and $\delta > 0$ there are $\varepsilon > 0$ and M such that, for any l -partite k -graph of order $n \geq M$ and size at least $t_l^k(n) - \varepsilon n^k$, the number of vertices in each part is

$$\left(\frac{1}{l} \pm \delta\right)n. \tag{2}$$

Proof of Theorem 4. Our proof uses induction on $k + l$. It is convenient to start with the trivial base case $l = k - 1$ which formally satisfies the conclusion of the theorem: F_k^k is the k -graph of one edge, and $T_{k-1}^k(n)$ has no edges. The other base case $k = 2$ is the content of the Simonovits stability theorem [21], so we further assume that $l \geq k > 2$.

Let $\delta = \delta_l > 0$ be given. Our goal is to obtain $\varepsilon = \varepsilon_l$ and $M = M_l$ satisfying the theorem. We choose the constants in this order:

$$\delta_l \gg \delta_{l-1} \gg \varepsilon_{l-1} \gg \varepsilon_l \gg \frac{1}{M_{l-1}} \gg \frac{1}{M_l},$$

where $a \gg b$ means that $b > 0$ is sufficiently small depending on a (and k, l). In particular, we assume that $\varepsilon_{l-1}, M_{l-1}$ demonstrate the validity of the theorem for $l - 1, k - 1$, and δ_{l-1} . Suppose that $n > M_l$. Let G be an \mathcal{F}_{l+1}^k -free k -graph on n vertices with

$$|G| \geq t_l^k(n) - \varepsilon_l n^k. \tag{3}$$

Pick a vertex $x \in V(G)$ of maximum degree Δ . Let $N = N(x)$ be the neighborhood of x , that is, the set of vertices $y \neq x$ for which $\text{codeg}_G(x, y) > 0$. Consider the k -graph $G[N]$ induced by N , and suppose that it contains a copy H of a member of \mathcal{F}_l^k . Let $C \subset V(H)$ be the core

of H , and $D \subset C$ for some $D \in G$. Form H' from H by adding the vertex x and edges containing each pair $\{x, v\}$ with $v \in C$. These edges exist by the definition of N . Therefore H' contains a member of \mathcal{F}_{l+1}^k with core $C \cup \{x\}$, which is a contradiction. Consequently, $G[N]$ is \mathcal{F}_l^k -free.

Next consider the $(k - 1)$ -graph L , where $L = L(x)$ is the link of x . Suppose that L contains a copy H of a member of \mathcal{F}_l^{k-1} . Enlarge every edge of H to contain x . The resulting k -graph contains a copy of some $H' \in \mathcal{F}_{l+1}^k$ with core $C \cup \{x\}$, a contradiction. Therefore L is \mathcal{F}_l^{k-1} -free.

Set $s = n - |N|$ and let $X = V(G) \setminus N$. Note that $x \in X$. Since $G[N]$ is \mathcal{F}_l^k -free and L is \mathcal{F}_l^{k-1} -free, Theorem 2 implies that $|G[N]| \leq t_{l-1}^k(n - s)$ and $\Delta = |L| \leq t_{l-1}^{k-1}(n - s)$. This gives

$$\begin{aligned} |G| &\leq |G[N]| + s \cdot \Delta - e_G(X) \\ &\leq t_{l-1}^k(n - s) + s \cdot t_{l-1}^{k-1}(n - s) - e_G(X) \\ &\leq t_l^k(n) - e_G(X), \end{aligned} \tag{4}$$

where the last inequality follows from (1). (Recall that $e_G(X)$ is the number of edges of G that intersect X in at least two vertices.) At this stage, one can deduce the upper bound in Theorem 2 by induction on $k + l$ since, obviously, $e_G(X) \geq 0$. (A further routine analysis will also show that $T_l^k(n)$ is the unique extremal configuration for $\text{ex}(n, \mathcal{F}_{l+1}^k)$.)

The inequalities (3) and (4) imply that

$$t_l^k(n) - \varepsilon_l n^k \leq t_{l-1}^k(n - s) + s \cdot t_{l-1}^{k-1}(n - s).$$

Note that the right-hand side is the size of the l -partite k -graph with n vertices such that one part has size s and the other $l - 1$ parts are almost equal. From (2), we conclude that

$$s = \left(\frac{1}{l} \pm \delta_{l-1} \right) n. \tag{5}$$

Moreover, routine calculations show (alternatively, see Claim 2 in [13, Theorem 3]) that (3) and (4) imply that

$$\Delta = |L| > t_{l-1}^{k-1}(n - s) - \varepsilon_{l-1}(n - s)^{k-1}. \tag{6}$$

Now consider L . This $(k - 1)$ -graph has $n - s$ vertices. Since $n \geq M_l \gg M_{l-1}$, we have $n - s \geq M_{l-1}$ by (5). Because of (6) we may apply the induction hypothesis to the \mathcal{F}_l^{k-1} -free $(k - 1)$ -graph L . We conclude that there exists a Turán hypergraph $T_{l-1} \cong T_{l-1}^{k-1}(n - s)$ with vertex partition $N = W_1 \cup \dots \cup W_{l-1}$ such that

$$|T_{l-1} \triangle L| \leq \delta_{l-1}(n - s)^{k-1}. \tag{7}$$

By (5) we conclude that for each $i \in [l - 1]$ we have

$$|W_i| = \frac{n - s}{l - 1} \pm 1 = \left(\frac{1}{l} \pm \delta_{l-1} \right) n. \tag{8}$$

Let $W_l = X$ and let T_l be the complete l -partite k -graph with the vertex partition $W_1 \cup \dots \cup W_l$. By (5) and (8) T_l is $\frac{\delta_l}{2} n^k$ -close to a $T_l^k(n)$ because we can transform one to the other by moving at most $\delta_{l-1} n \times l$ vertices between parts, thus changing at most $\delta_{l-1} l n \times \binom{n-1}{k-1} < (\delta_l/2) n^k$ edges.

We will show that

$$|G \setminus T_l| \leq \frac{\delta_l}{5} n^k. \tag{9}$$

This implies, in view of (3) and the inequality $|T_l| \leq t_l^k(n)$, that

$$|G \Delta T_l| = |T_l| - |G| + 2|G \setminus T_l| \leq \varepsilon_l n^k + \frac{2\delta_l}{5} n^k < \frac{\delta_l}{2} n^k,$$

and the desired bound $|G \Delta T_l^k(n)| \leq \delta_l n^k$ follows from the triangle inequality.

From (4) we conclude that $e_G(X) \leq \varepsilon_l n^k$. Suppose on the contrary to (9) that we have more than $\frac{\delta_l}{5} n^k - \varepsilon_l n^k > \frac{\delta_l}{6} n^k$ edges of G intersecting some part of N in at least two vertices. By averaging there is an $i \in [l - 1]$ such that $|B| \geq \frac{\delta_l}{6l} n^2$, where B consists of all 2-subsets of W_i covered by at least one edge of G . Assume that $i = l - 1$ without loss of generality.

Let $w = (\frac{1}{7} - \delta_{l-1})n$. Recall that w is a lower bound on each $|W_i|$ by (5) and (8). For every choice of $x_1 \in W_1, \dots, x_{l-2} \in W_{l-2}$ and $\{x_{l-1}, x_l\} \in B$, at least $w^{l-2} \times \frac{\delta_l}{6l} n^2$ choices in total, we consider a potential copy of \mathcal{F}_{l+1}^k with core $C = \{x, x_1, \dots, x_l\}$. (Recall that x is the chosen vertex of maximum degree.) As G is \mathcal{F}_{l+1}^k -free, at least one of the following must hold:

1. $K \notin G$, where $K = \{x, x_1, \dots, x_{k-1}\}$.
2. A pair $\{x, x_i\}$ with $i \in [l]$ is not covered by an edge of G .
3. A pair $\{x_i, x_j\}$ with $\{i, j\} \neq \{l - 1, l\}$ is not covered by an edge of G .

One of these three alternatives holds for at least one third of the choices of x_i 's. If it is Alternative 1, then for each such K we have $K \setminus \{x\} \in T_{l-1} \setminus L$. Any fixed set K is counted at most n^{l-k+1} times. Now, since $\delta_{l-1} \ll \delta_l$, we obtain a contradiction to (7):

$$|T_{l-1} \setminus L| \geq \frac{1}{3} \times w^{l-2} \times \frac{\delta_l}{6l} n^2 \times n^{-l+k-1} > \delta_{l-1} (n - s)^{k-1}.$$

If it is Alternative 2, then we obtain a contradiction as follows. For every uncovered pair $\{x, x_i\}$, the vertex x_i belongs to at least w^{k-2} edges of the $(k - 1)$ -graph T_{l-1} . None of these edges belongs to L , for otherwise the pair $\{x, x_i\}$ would be covered by an edge of G . On the other hand, every edge $D \in T_{l-1} \setminus L$ appears this way for at most $(k - 1)n^{l-1}$ choices of the sequence (x_1, \dots, x_l) : we have to choose $x_i \in D$ and then the other $l - 1$ vertices x_j . Thus we have

$$|T_{l-1} \setminus L| \geq \frac{1}{3} \times w^{l-2} \times \frac{\delta_l}{6l} n^2 \times w^{k-2} \times \frac{1}{(k - 1)n^{l-1}} > \delta_{l-1} (n - s)^{k-1},$$

again a contradiction to (7). Finally, suppose that Alternative 3 appears frequently. Each pair $\{x_i, x_j\}$ belongs to at least w^{k-3} edges of $T_{l-1} \setminus L$. However, each such edge is counted at most $\binom{k-1}{2} n^{l-2}$ times. Hence,

$$|T_{l-1} \setminus L| \geq \frac{1}{3} \times w^{l-2} \times \frac{\delta_l}{6l} n^2 \times w^{k-3} \times \binom{k - 1}{2}^{-1} n^{-l+2} > \delta_{l-1} (n - s)^{k-1}.$$

Again we obtain a contradiction to (7). This completes the proof of Theorem 4. \square

3. Step 3: F_{l+1}^k is stable

Please note that Theorem 5 below is formally stronger than Theorem 4. However, it follows from Theorem 4 by an application of Lemma 4 from [15]. The last result indirectly relies on the recent Hypergraph Regularity Lemma of Gowers [7] or Nagle–Rödl–Schacht–Skokan [14, 17, 18]. For our particular hypergraph F_{l+1}^k , the recourse to such a complicated technique is not

really necessary and we present a short and self-contained proof, similar to the proof of Lemma 3 in [15].

Theorem 5. For any $l \geq k \geq 2$ and $\delta > 0$ there exist $\varepsilon > 0$ and M such that the following holds for all $n > M$: Any n -vertex F_{l+1}^k -free k -graph G with at least $t_l^k(n) - \varepsilon n^k$ edges is δn^k -close to $T_l^k(n)$.

Proof. Given $\delta > 0$, let $\delta \gg \varepsilon \gg 1/M$.

Suppose that $n > M$ and G is an n -vertex F_{l+1}^k -free k -graph with at least $t_l^k(n) - \varepsilon n^k$ edges. Let G' be obtained from G by deleting all edges that contain a pair of vertices whose codegree is at most $l^3 \binom{n}{k-3}$. Since the number of pairs of vertices is $\binom{n}{2}$, we have

$$|G \setminus G'| \leq l^3 \binom{n}{k-3} \times \binom{n}{2} < \varepsilon n^k < \frac{\delta}{2} n^k. \tag{10}$$

Now we argue that G' is \mathcal{F}_{l+1}^k -free. Suppose on the contrary that G' contains a copy of some $F \in \mathcal{F}_{l+1}^k$ with core C and edge $D \subset C$. Since every pair of vertices $x, y \in C$ is contained in an edge of G' , we have, by $l \geq k \geq 2$,

$$\text{codeg}_G(x, y) \geq l^3 \binom{n}{k-3} > \left(\binom{l+1}{2} (k-2) + l + 1 \right) \binom{n}{k-3}.$$

Hence we can greedily choose edges of G containing all pairs in $\binom{C}{2} \setminus \binom{D}{2}$, so that these edges intersect C in precisely two vertices and are pairwise disjoint outside C . The resulting set of $\binom{l+1}{2} - \binom{k}{2}$ edges, together with D , forms a copy of F_{l+1}^k in G , a contradiction.

We have

$$|G'| > |G| - \varepsilon n^k \geq (t_l^k(n) - \varepsilon n^k) - \varepsilon n^k = t_l^k(n) - 2\varepsilon n^k.$$

We apply Theorem 4 to G' and conclude that G' is $\frac{\delta}{2} n^k$ -close to $T_l^k(n)$. By (10), G and $T_l^k(n)$ are δn^k -close. The proof is complete. \square

4. Step 4: Proof of Theorem 3

If $k = 2$, then Theorem 3 is precisely the Turán theorem [22]. Thus let us assume that $l \geq k \geq 3$. Choose small $c \gg c' \gg \delta > 0$. Let n be large.

Let G be an F_{l+1}^k -free k -graph on $[n]$ with $|G| = t_l^k(n)$. We will show that G is l -partite. This implies the theorem because $T_l^k(n)$ is the unique l -partite k -graph on n vertices with $t_l^k(n)$ edges, and the addition of any edge to $T_l^k(n)$ yields a copy of F_{l+1}^k .

Let $W_1 \cup \dots \cup W_l$ be a partition of $[n]$ such that

$$f = \sum_{D \in G} |\{i \in [l]: D \cap W_i \neq \emptyset\}|$$

is maximum possible. Let T be the complete l -partite k -graph on $W_1 \cup \dots \cup W_l$. Let us call the edges in $T \setminus G$ missing and the edges in $G \setminus T$ bad. As $|T| \leq t_l^k(n) = |G|$, the number of bad edges is at least the number of missing edges.

By Theorem 5, there is an l -partite k -graph which is δn^k -close to G . Consequently, $f \geq k(|G| - \delta n^k)$. On the other hand,

$$f \leq k|G \cap T| + (k-1)|G \setminus T| = k|G| - |G \setminus T|.$$

This implies that $|G \setminus T| \leq k\delta n^k$ and, in view of $|T| \leq |G|$,

$$|T \setminus G| \leq k\delta n^k. \tag{11}$$

Thus we have $|T| \geq |G \cap T| \geq t_l^k(n) - k\delta n^k$. From (2) we conclude that for each $i \in [l]$ we have, for example, $||W_i| - \frac{n}{7}| \leq \frac{n}{27}$.

If $G \subset T$, then we are done. Thus, let us assume that B is non-empty, where the 2-graph B consists of all *bad* pairs, that is, pairs of vertices which come from the same part W_i and are covered by an edge of G .

For distinct vertices x, y call the pair $\{x, y\}$ *sparse* if G has at most $\binom{l+1}{2}(k-2) + l + 1 \binom{n}{k-3}$ edges containing both x and y ; otherwise $\{x, y\}$ is called *dense*. It is easy to see that if we have a fixed $(l+1)$ -set $C \subset V(G)$ containing at least one edge $D \in G$, then at least one pair of vertices $\{x, y\}$ from $\binom{C}{2} \setminus \binom{D}{2}$ is sparse. (For otherwise we can greedily build a copy of F_{l+1}^k in G with the core C .)

Let A consist of those $z \in V(G)$ which are incident to at least cn^{k-1} missing edges.

Claim 1. Any bad pair $\{x_0, x_1\}$ intersects A .

Proof. Assume without loss of generality that $x_0, x_1 \in W_1$ are covered by $D \in G$.

It is easy to see that for any choice of $x_i \in W_i \setminus D$ for $i \in [2, l]$ (at least $(\frac{n}{27} - k)^{l-1} > (\frac{n}{37})^{l-1}$ choices), at least one pair $\{x_i, x_j\}$ with $\{i, j\} \neq \{0, 1\}$ is sparse or the k -tuple $\{x_1, x_2, \dots, x_k\}$ is missing for otherwise we obtain a copy of F_{l+1}^k . (In fact, we can make stronger claims but this one suffices.)

If the second alternative occurs at least a half of the time, then $x_1 \in A$. Indeed, any k -tuple $D \ni x_1$ is counted at most n^{l-k} times (the number of ways to choose x_{k+1}, \dots, x_l), so x_1 belongs to at least $\frac{1}{2}(\frac{n}{37})^{l-1}/n^{l-k} \geq cn^{k-1}$ missing edges, as required.

So, suppose that for at least half of the choices, the first alternative holds, i.e., there is a sparse pair. Each such pair $\{x_i, x_j\}$ appears, very roughly, at most n^{l-3} times unless $\{x_i, x_j\} \cap \{x_0, x_1\} \neq \emptyset$ when the pair is counted at most n^{l-2} times. There are two further alternatives to consider.

If at least a quarter of the time, the found sparse pair is disjoint from $\{x_0, x_1\}$, then we obtain at least $\frac{1}{4}(\frac{n}{37})^{l-1}/n^{l-3} \geq cn^2$ sparse pairs, each intersecting two parts W_i . But this leads to a contradiction to (11): each such sparse pair is contained in at least, say, $(\frac{n}{37})^{k-2}$ missing edges while each missing edge contains at most $\binom{k}{2}$ sparse pairs. Hence, at least a quarter of the time, the sparse pair intersects $\{x_0, x_1\}$, so one of these vertices, say x_0 , is in at least $\frac{1}{8}(\frac{n}{37})^{l-1}/n^{l-2}$ sparse pairs, which implies that $x_0 \in A$. The claim has been proved. \square

Considering vertices from A , we obtain at least $|A| \times cn^{k-1}/k$ missing edges and, consequently, at least $|A| \times cn^{k-1}/k$ bad edges. Let \mathcal{B} consist of the pairs $(D, \{x, y\})$, where $\{x, y\} \in B$, $D \in G$ and $x, y \in D$. (Thus D is a bad edge.) As each bad edge contains at least one bad pair, we conclude that $|\mathcal{B}| \geq |A| \times cn^{k-1}/k$. For any $(D, \{x, y\}) \in \mathcal{B}$, we have $\{x, y\} \cap A \neq \emptyset$ by Claim 1. If we fix x and D , then, obviously, there are at most $k - 1$ ways to choose a bad pair $\{x, y\} \subset D$. By Claim 1, some vertex $x \in A$, say $x \in W_1$, belongs to at least

$$\frac{|\mathcal{B}|}{(k-1)|A|} \geq \frac{c}{k(k-1)}n^{k-1}$$

bad edges, each intersecting W_1 in another vertex y .

Let $Y_1 \subset W_1$ be the neighborhood of x in the 2-graph B . Let $Z_1 \subset Y_1$ be the set of those vertices z for which $\{x, z\}$ is dense. The number of edges containing x and some vertex of $Y_1 \setminus Z_1$ is at most $l^3 n^{k-2} < \frac{c}{2k(k-1)} n^{k-1}$. Consequently, the number of bad edges containing x and some vertex of Z_1 is at least $\frac{c}{2k(k-1)} n^{k-1}$. Therefore $|Z_1| \geq \frac{c}{2k(k-1)} n \geq c'n$.

Let Z_j consist of those $z \in W_j$ for which $\{x, z\}$ is dense, $j \in [2, l]$. If $|Z_j| \geq c'n$ for each $j \in [2, l]$, then every l -tuple (x_1, x_2, \dots, x_l) with $x_j \in Z_j$ (at least $(c'n)^l$ choices) generates a sparse pair not containing x or the edge $\{x_1, \dots, x_k\}$ is missing. The latter alternative cannot happen, say, at least half of the time because otherwise we obtain more than $\frac{1}{2}(c'n)^l/n^{l-k} > k\delta n^k$ missing edges, a contradiction to (11). Thus at least half of the time, we obtain a sparse pair disjoint from x . This gives at least $\frac{1}{2}(c'n)^l/n^{l-2}$ sparse pairs, each intersecting some two parts, which leads to a contradiction to (11).

Hence, assume that, for example, $|Z_2| < c'n$. This means that all but at most $c'n$ pairs $\{x, z\}$ with $z \in W_2$ are sparse, that is, there are at most $l^3 n^{k-2} + c'n^{k-1} < 2c'n^{k-1}$ G -edges containing x and intersecting W_2 . Let us contemplate moving x from W_1 to W_2 . Some edges of G may decrease their contribution to f . But each such edge must contain x and intersect W_2 so the corresponding decrease is at most $2c'n^{k-1}$. On the other hand, the number of edges of G containing x , intersecting $W_1 \setminus \{x\}$, and disjoint from W_2 is at least $(\frac{c}{k(k-1)} - 2c')n^{k-1} > 2c'n^{k-1}$. Hence, by moving x from W_1 to W_2 we strictly increase f , a contradiction to the choice of the parts W_i . The theorem is proved.

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