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European Journal of Combinatorics

European Journal of Combinatorics 26 (2005) 607-616

www.elsevier.com/locate/ejc

# Integer sets with prescribed pairwise differences being distinct

Béla Bollobás<sup>a,b</sup>, Oleg Pikhurko<sup>c</sup>

<sup>a</sup>University of Memphis, Memphis, TN 38152-3240, USA <sup>b</sup>Trinity College, Cambridge CB2 1TQ, UK <sup>c</sup>Department of Mathematical Sciences, Carnegie Mellon University, Pittsburgh, PA 15213, USA

Received 20 January 2004; received in revised form 25 April 2004; accepted 26 April 2004 Available online 10 June 2004

#### Abstract

We label the vertices of a given graph G with positive integers so that the pairwise differences over its edges are all distinct. Let  $\mathcal{D}(G)$  be the smallest value that the largest label can have.

For example, for the complete graph  $K_n$ , the labels must form a Sidon set. Hence,  $\mathcal{D}(K_n) = (1 + o(1))n^2$ . Rather surprisingly, we demonstrate that there are graphs with only  $n^{\frac{3}{2}+o(1)}$  edges achieving this bound.

More generally, we study the maximum value of  $\mathcal{D}(G)$  that a graph G of the given order n and size m can have. We obtain bounds which are sharp up to a logarithmic multiplicative factor. The analogous problem for pairwise sums is considered as well. Our results, in particular, disprove a conjecture of Wood.

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# 1. Introduction

Let G be a graph. A *difference-magic labelling* of G is an injective mapping  $l: V(G) \rightarrow \mathbb{N}$  (into positive integers) such that the e(G) numbers

 $|l(x) - l(y)|, \quad \{x, y\} \in E(G),$ 

are pairwise distinct.

It is trivial to see that every graph admits a difference-magic labelling, so a natural question to ask is how economical it can be. More precisely, we should like to determine

E-mail address: bollobas@msci.memphis.edu (B. Bollobás).

URL: http://www.math.cmu.edu/~pikhurko/ (O. Pikhurko).

the *difference-magic number*  $\mathcal{D}(G)$  which is the smallest k such that a difference-magic labelling of G into  $[k] := \{1, \ldots, k\}$  exists.

For example, it is easy to see that if *G* is the complete graph of order *n*, then  $\mathcal{D}(G)$  is precisely  $s_n$ , the smallest *s* such that [s] contains a *Sidon* subset of size *n*. (A set  $A \subset \mathbb{Z}$  is *Sidon* if all sums a + b with  $a, b \in A$  and  $a \leq b$  are distinct.) The latter problem is well studied; the results of Singer [13] and Erdős and Turán [8] (see e.g. Halberstam and Roth [10, Chapter II]) imply that  $s_n = (1 + o(1))n^2$ . Erdős [5] offered \$500 for proving or disproving that  $s_n = n^2 + O(n)$ .

Here we deal with

$$\mathcal{D}(n,m) := \max\{\mathcal{D}(G) : v(G) = n, \ e(G) = m\},\$$

the maximum value of  $\mathcal{D}(G)$  for a graph G of order n and size m.

It turns out that

$$\mathcal{D}(n,m) = (1+o(1))n^2 \qquad \text{if } m / \sqrt{n^3 \ln n} \to \infty .$$
<sup>(1)</sup>

In fact, a random graph of order *n* with the appropriate edge probability demonstrates (1). We find it surprising that graphs so sparse (with only  $n^{\frac{3}{2}+o(1)}$  edges) have the  $\mathcal{D}$ -function asymptotically the same as that of the complete graph.

What happens for smaller *m*? The obvious choice is to consider random graphs of suitable density. This, indeed, leads to interesting results. Let  $G \in \mathcal{G}(n, p)$ , that is, G is a random graph on *n* vertices where each edge is included in G independently of others and with probability *p*. If  $p = O((\ln n/n)^{1/2})$  and  $p > n^{-1+\varepsilon}$ , then

$$\mathcal{D}(G) = \Theta(n^3 p^2 / \ln n). \tag{2}$$

A lower bound on  $\mathcal{D}(n, m)$  can be obtained by adding isolated vertices to a random graph and figuring out the best parameters to choose. On the other hand, the simple labelling procedure described in Section 4 gives an upper bound that is within an  $O((\ln n)^{2/3})$ -factor of the lower bound. Roughly, we obtain

$$\mathcal{D}(n,m) = m^{4/3 + O\left(\frac{\ln \ln m}{\ln m}\right)}, \qquad \text{if } m = O\left(\sqrt{n^3 \ln n}\right), \tag{3}$$

unless  $m = o(n^{3/4})$  when  $\mathcal{D}(n,m) = (1 + o(1))n$ . All details (with more precise expressions for the error terms) can be found in the corresponding sections.

Let us define a *sum-magic labelling* of a graph G as an injection  $l : V(G) \to \mathbb{N}$  such that all e(G) sums l(x) + l(y),  $\{x, y\} \in E(G)$ , are pairwise distinct. We ask for the *sum-magic number* S(G), the smallest value that the largest label can have, and for

$$\mathcal{S}(n,m) := \max\{\mathcal{S}(G) : v(G) = n, e(G) = m\}.$$

It is not surprising that most of the methods on the  $\mathcal{D}$ -function transfer to S, giving similar bounds. (In particular, (3) holds for S(n, m) as well.) However, there is one peculiar distinction. While Corollary 2 states that  $S(K_n) = (1 + o(1))n^2$ , Theorem 3 shows that there is a constant c > 0 such that  $S(n, m) < (1 - c)n^2$  whenever  $m \le cn^2$ . Random graphs are far worse in hitting  $(1 + o(1))n^2$ : this happens only when the random graph is almost complete.

Wood [15] defines an *edge-magic injection* with the *magic sum s* as an injection  $l : V(G) \cup E(G) \rightarrow \mathbb{N}$  such that for any edge  $\{a, b\} \in E(G)$  we have  $s = l(a) + l(b) + l(\{a, b\})$ . Let  $\mathcal{E}(G)$  be the smallest possible value of *s*. Wood [15, Section 7] conjectured that there is an absolute constant *C* such that for any graph *G* we have  $\mathcal{E}(G) \leq C(v(G) + e(G))$ . Clearly, the vertex labels of any edge-magic injection form a sum-magic labelling, so  $\mathcal{E}(G) \geq \mathcal{S}(G)$  and random graphs disprove Wood's conjecture.

One can also ask what is the value of, for example,

$$\mathcal{S}_{\min}(n,m) := \min\{\mathcal{S}(G) : v(G) = n, e(G) = m\}.$$

This is the inverse problem to maximising the number of distinct pairwise sums that a set  $A \subset [s]$  of given size *n* can have. This question is investigated by Pikhurko [12].

# 2. Some preliminary results

Let  $A \in {[m] \choose n}$ , meaning that A is an *n*-subset of [m].

Recall that A is called a *Sidon set* if the sums a + b,  $a, b \in A$  with  $a \ge b$ , are pairwise distinct, which is equivalent to all differences a - b,  $a, b \in A$  with a > b, being pairwise distinct. Erdős and Turán [8] proved that this property implies that  $m \ge (1 + o(1))n^2$ . The following results show that, in a sense, it is the condition on differences (rather than that on sums) which pushes max A upwards.

For  $i \in [m-1]$  let  $g_i$  be the number of representations i = a - b with  $a, b \in A$ . Thus, if  $n^2 \ge (1 + \varepsilon)m$  (and *n* is large), then there must be *i* with  $g_i \ge 2$ . Although the following theorem strengthens this claim considerably, its proof goes via an easy modification of the original argument of Erdős and Turán [8]. A similar result (in a more precise form) was independently obtained by Ferrara, Kohayakawa and Rödl [9, Lemma 12].

Let  $f_+ = f$  if f > 0 and  $f_+ = 0$  otherwise.

**Theorem 1.** Let  $\varepsilon > 0$  be fixed and  $n \to \infty$ . Let  $A \in {\binom{[m]}{n}}$ . If  $n^2 \ge (1 + \varepsilon)m$ , then  $g = \Omega(n^2)$ , where  $g := \sum_{i=1}^{m-1} (g_i - 1)_+$ .

**Proof.** Let  $t := cn^2$ , where  $c = c(\varepsilon) > 0$  is a small constant. Assume  $t \in \mathbb{N}$ . Define

$$A_i := A \cap [i, i + t - 1]$$
 and  $a_i := |A_i|, \quad i \in [2 - t, m],$ 

where  $[i, j] := \{i, i + 1, ..., j\}.$ 

Let  $\mathcal{X}$  consist of all quadruples (a, b, i, x) such that x = a - b > 0 and  $a, b \in A_i$ . Using the identity  $\sum_{i=2-t}^{m} a_i = nt$  and the quadratic-arithmetic mean inequality, we obtain

$$|\mathcal{X}| = \sum_{i=2-t}^{m} \binom{a_i}{2} = \frac{1}{2} \sum_{i=2-t}^{m} a_i^2 - \frac{nt}{2} \ge \frac{(nt)^2}{2(m+t-1)} - \frac{nt}{2}.$$
(4)

Each  $x \in [t-1]$  is included in  $g_x \cdot (t-x) \le (t-x) + t(g_x - 1)_+$  quadruples. Hence,

$$|\mathcal{X}| \le \sum_{x=1}^{t-1} (t - x + t(g_x - 1)_+) = \frac{t(t-1)}{2} + gt.$$
(5)

By choosing *c* sufficiently small, we can ensure that the right-hand side of (4) is, for example, at least  $(1 + \frac{\varepsilon}{2})\frac{t^2}{2}$ , which together with (5) implies the theorem.  $\Box$ 

We will need Theorem 1 in Section 3. Here we demonstrate another application.

**Corollary 2.**  $S(K_n) = (1 + o(1))n^2$ .

**Proof.** Let *A* be the label set of a sum-magic labelling. Note that *A* need not be Sidon as it may well happen that a - c = c - b for  $a, b, c \in A$ . However, if a - b = c - d with  $a \notin \{b, c\}$ , then either a = d or b = c. It follows that  $g_x \le 2$  for any x > 0 and, if  $g_x = 2$ , then there are  $a, b, c \in A$  with a - b = c - a = x. If  $a - b' = c' - a \neq 0$ , then we have b' + c' = 2a = b + c and thus  $\{b', c'\} = \{b, c\}$ . Hence, no *a* can appear for more than one *x* in the above manner. We conclude that  $g \le |A|$ , implying the claim by Theorem 1.  $\Box$ 

The natural analogue of Theorem 1 in terms of the number of solutions to x = a + b,  $a, b \in A$ , is not true, as the following construction of Erdős and Freud [6] demonstrates. Let  $S \in {\binom{[t]}{s}}$  be a Sidon set with  $t = (1 + o(1))s^2$ . (Such sets were constructed by Singer [13].) Let  $X = S \cup S'$ , where

$$S' := 3t + 1 - S := \{3t + 1 - a : a \in S\} \subset [2t + 1, 3t].$$

Clearly,  $S + S \subset [2, 2t]$ ,  $S + S' \subset [2t+2, 4t]$  and  $S' + S' \subset [4t+2, 6t]$  are disjoint. Hence, all sums a + b,  $a, b \in X$  with  $a \le b$ , are pairwise distinct except those *s* sums which are equal to 3t + 1. If the complement of an order-*n* graph *G* has a matching covering all but r = o(n) vertices, then considering the first *n* elements of the set *X* constructed above for  $s := \frac{n+r}{2}$ , we conclude

$$S(G) \le (3/4 + o(1))n^2.$$
 (6)

By modifying the above construction, we can show one of the results claimed in the Introduction.

**Theorem 3.** There is a constant c > 0 such that if  $m \le cn^2$ , then

$$\mathcal{S}(n,m) \le (1-c)n^2. \tag{7}$$

**Proof.** Let  $\alpha = 0.9$ , for example. In the above construction of  $X = S \cup S'$  let  $Y \subset X$  consist of the first  $n := \lfloor (1 + \alpha)s \rfloor$  elements of X. As it was shown by Erdős and Freud [6, Lemma 1], any asymptotically maximum Sidon subset of [t] is almost uniformly distributed. This implies that max  $Y = (2 + \alpha + o(1))t$ .

Now, all sums in Y + Y are distinct except those sums which equal 3t + 1. The number of these exceptional sums is  $\lfloor \alpha s \rfloor = (\frac{\alpha}{1+\alpha} + o(1))n$ . So, if the complement of an order-*n* graph *G* has a matching of size bigger than  $0.48n > (\frac{\alpha}{1+\alpha} + o(1))n$ , then

$$S(G) \le (2 + \alpha + o(1))t = \frac{2 + \alpha + o(1)}{(1 + \alpha)^2}n^2 < 0.9n^2.$$

It follows from the Tutte 1-factor theorem [14] that a matching of size 0.48*n* in the complement  $\overline{G}$  is guaranteed if  $e(G) \leq \delta n^2$  for some constant  $\delta > 0$ . Now, the theorem follows.  $\Box$ 

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**Remark.** Random graphs do not provide good examples if we want to achieve  $S(G) = (1+o(1))n^2$ : this happens only when  $1-p = O(\frac{\ln n}{n})$ . Indeed, Erdős and Rényi [7] (cf. Bollobás and Thomason [3]) showed that if  $p \le 1 - (1+\varepsilon)\frac{\ln n}{2n}$  then with high probability the complement of  $G \in \mathcal{G}(n, p)$  has an almost perfect matching; so then (6) holds.

#### 3. Random graphs

**Theorem 4.** Fix any  $\delta > 0$ . Let  $G \in \mathcal{G}(n, p)$ , where  $n \to \infty$  and  $p \in (0, 1)$  is a function of n such that  $np/\ln n \to \infty$ . Let  $\lambda := p\sqrt{n/\ln n}$ . Then almost surely  $\mathcal{D}(G) \ge d$  and  $\mathcal{S}(G) \ge s$ , where

$$d := \begin{cases} (1-\delta)n^2 \\ \left(\frac{\lambda^2}{16+2\lambda^2}-\delta\right)n^2, \\ \left(\frac{1}{16}-\delta\right)\frac{n^3p^2}{\ln n}, \end{cases} \quad s := \begin{cases} \left(\frac{1}{4}+\frac{1}{(\pi+2)^2}-\delta\right)n^2, & \text{if } \lambda \to \infty, \\ \left(\frac{\lambda^2}{32+4\lambda^2}-\delta\right)n^2, & \text{if } \lambda = \Theta(1), \\ \left(\frac{1}{32}-\delta\right)\frac{n^3p^2}{\ln n}, & \text{if } \lambda = o(1). \end{cases}$$

**Proof.** We prove the lower bound on  $\mathcal{D}(G)$ . Let [n] be the vertex set. Let  $\varepsilon > 0$  be a small constant depending on  $\delta$ . Assume  $d \in \mathbb{N}$ .

Fix an injective mapping  $l : [n] \to [d]$ . Now, let us choose  $G \in \mathcal{G}(n, p)$ . We want to bound the probability p' that all differences l(i) - l(j), with  $\{i, j\} \in {[n] \choose 2}$  being an edge of  $\mathcal{G}(n, p)$  and l(i) > l(j), are pairwise distinct. If u is an upper bound on p' for any l, then the probability that  $G \in \mathcal{G}(n, p)$  satisfies  $\mathcal{D}(G) \le d$  is at most  $n! {d \choose n} u < d^n u$ . Hence, if we can show that  $p' = o(d^{-n})$ , then almost surely  $\mathcal{D}(G) > d$ .

For  $k \in [d]$ , let  $g_k$  be the number of representations k = l(i) - l(j) with  $i, j \in [n]$ . Let  $t := \binom{n}{2} = \sum_{k=1}^{d} g_k$ . Clearly,

$$p' = \prod_{k=1}^{d} p_k,\tag{8}$$

where  $p_k = (1-p)^{g_k} + g_k p(1-p)^{g_k-1}$  is the probability of selecting at most one edge with difference k. (Note that the formula is also valid for  $g_k = 0$  and  $g_k = 1$ , when  $p_k = 1$ .) It is routine to see that

$$p' = \prod_{k=1}^{d} p_k \le ((1-p)^{t/d} + (pt/d)(1-p)^{(t/d)-1})^d.$$
(9)

**Case 1.**  $p = o(\sqrt{\ln n/n})$ , that is,  $\lambda = o(1)$ .

We have  $t/d \to \infty$  and pt/d = o(1). Using the inequality  $e^{-x}(1+x) \le 1 - (\frac{1}{2} - \varepsilon)x^2$  valid if x > 0 is small, we deduce from (9) the required bound on p':

$$p' \le ((1-p)^{t/d}(1+pt/d+2p^2t/d))^d \le (e^{-pt/d}(1+pt/d+2p^2t/d))^d$$
  
$$\le (1-(1/2-\varepsilon)(pt/d)^2+2p^2t/d)^d \le e^{-(1/2-2\varepsilon)(pt)^2/d} = o(e^{-n\ln d}).$$

**Case 2.**  $p = \Theta(\sqrt{\ln n/n})$ , that is,  $\lambda = \Theta(1)$ .

We have t/d = O(1) so we can simply take the Taylor expansion of (9) to obtain the required bound:

$$p' \le \left(1 + \left(\frac{t}{2d} - \frac{t^2}{2d^2}\right)p^2 + O(p^3)\right)^d \le e^{\frac{1}{2}(t - \frac{t^2}{d} + \varepsilon)p^2} = o(e^{-n\ln d}).$$

**Case 3.**  $p\sqrt{n/\ln n} \to \infty$ , that is,  $\lambda \to \infty$ .

By Theorem 1 we know that  $g := \sum_{k=1}^{d} (g_k - 1)_+ = \Omega(n^2)$ . It is routine to see that if  $g_i \ge g_j + 2$ , then the right-hand side of (8) increases if we replace  $g_i$  and  $g_j$  by  $g_i - 1$  and  $g_j + 1$  respectively. Hence,

$$p' \le ((1-p)^2 + 2p(1-p))^g = (1-p^2)^g = o(d^{-n}),$$

as required.

Let us turn to the sum-magic number. Fix an injection  $l : [n] \to [s]$ . For  $k \in [2s]$  define  $g_k$  as the number of representations k = l(i) + l(j) with  $1 \le i < j \le n$ . Let  $t := \binom{n}{2} = \sum_{k=1}^{2s} g_k$ . The remainder of the proof goes via the obvious modification of the argument for  $\mathcal{D}(G)$  except that for  $\lambda \to \infty$  we use the result by Pikhurko [12, Theorem 2] which implies that  $\sum_{k=1}^{2s} (g_k - 1)_+ = \Omega(n^2)$ . (If we are content with  $s = (\frac{1}{4} - \delta)n^2$ , then  $\sum_{k=1}^{2s} (g_k - 1)_+ = \Omega(n^2)$  follows by trivial counting.) The reader should have little difficulty in filling in all missing details.  $\Box$ 

**Remark.** There is a jump in the lower bounds when we change from the case  $\lambda = \Theta(1)$  to  $\lambda \to \infty$ . It should be possible to 'smoothen' this by improving our bounds for large but bounded  $\lambda$ . However, the calculations seem to be rather unpleasant, so we do not go into the details.

**Remark.** As it was mentioned in the introduction, Theorem 4 disproves the conjecture of Wood in view of the inequality  $\mathcal{E}(G) \geq \mathcal{S}(G)$ . Indeed, if we take  $G \in \mathcal{G}(n, n^{-1/2})$  for example, then almost surely  $e(G) = (\frac{1}{2} + o(1))n^{3/2}$  while  $\mathcal{D}(G) = \Omega(n^2/\ln n)$ . With a bit of extra work it is possible to show that under the assumptions of Theorem 4 we have almost surely  $\mathcal{E}(G) \geq 2s$ . To do this, prove that, almost surely, any sum-magic labelling of *G* has  $\Omega(n)$  labels which are greater than *s* and there is an edge connecting two such labels. We leave the details to the interested reader.

Now let us turn to upper bounds.

**Theorem 5.** Let  $\delta > 0$  be fixed. Let  $G \in \mathcal{G}(n, p)$ , where  $n \to \infty$  and  $p \in (0, 1)$  is a function of n such that  $\frac{\ln(np)}{\ln \ln n} \to \infty$  and  $p = O(((\ln n/n)^{1/2}))$ . Then almost surely  $\mathcal{D}(G) \leq 2m$  and  $\mathcal{S}(G) \leq m$ , where

$$m \coloneqq (1+\delta)\frac{n^3 p^2}{\ln(np)}.$$
(10)

**Proof.** Let us estimate  $\mathcal{S}(G)$ . (The case of  $\mathcal{D}(G)$  is dealt with almost identically.)

We can assume that  $\delta$  is sufficiently small and  $m \in \mathbb{N}$ . Let *n* be large and  $\varepsilon > 0$  be a small constant depending on  $\delta$ . Let V(G) = [n] be the vertex set. Chernoff's bound [4]

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implies that almost surely we have

$$\left|\left|\Gamma(i+1)\cap[i]\right| - ip\right| \le \varepsilon np, \quad \text{for all } i \in [0, n-1], \tag{11}$$

where  $\Gamma(i + 1)$  is the set of neighbours of  $i + 1 \in V(G)$ .

Consider the conditional distribution of G given (11). We have gained the very useful control over the edges while some important properties of  $G \in \mathcal{G}(n, p)$  are preserved. (That is, almost sure events stay so; the random set  $\Gamma(i + 1) \cap [i]$  is independent from G[[i]], etc.)

We choose vertex labels one by one, doing the label arithmetic in  $M = \mathbb{Z}/m\mathbb{Z}$  (that is, modulo *m*). Our labelling  $l : V(G) \to [m]$  will have the property that the sums l(x) + l(y),  $\{x, y\} \in E(G)$ , will be pairwise distinct modulo *m*.

Suppose that we have already chosen labels for the vertices in I := [i].

Let

$$K := \{l(x) + l(y) : \{x, y\} \in E(G), x, y \in I\} \subset M$$

and k := |K|. By (11),

$$k \le \left(\frac{1}{2} + \varepsilon\right) inp. \tag{12}$$

Clearly, we can find a suitable label for i + 1 if

$$M \setminus l(I) \not\subset \bigcup_{x \in I \cap \Gamma(i+1)} (K - l(x)), \tag{13}$$

that is, if the translates  $K - l(x), x \in I \cap \Gamma(i + 1)$ , do not cover  $M \setminus l(I)$ .

This is obviously the case if

$$|M \setminus (\bigcup_{x \in I} (K - l(x)))| \ge n,$$

so let us assume otherwise. Then we have  $m - ik \le n$ , which implies by (12) that

$$i \ge n\sqrt{2p/\ln(np)}.\tag{14}$$

Now, we have to overcome the difficulty that *i* is large enough to potentially refute (13). In outline, we fix the labelling *l* of *I* and then choose the random set  $I \cap \Gamma(i + 1)$ . The labels  $l(x), x \in I \cap \Gamma(i + 1)$ , are random variables. If the translates K - l(x) cover the whole of  $M \setminus l(I)$ , then for every  $z \in M \setminus l(I)$  at least one element  $l(x) \in K - z$  is chosen. We prove that this is unlikely.

Let *S* consist of those elements from  $M \setminus l(I)$  which are covered by at most

$$t := \lfloor (1 + \varepsilon) kn/m \rfloor$$

of the translates  $K - l(x), x \in I$ . Clearly,

$$|M \setminus (S \cup l(I))| \times (1 + \varepsilon) kn/m \le kn,$$

so

$$s := |S| \ge \frac{\varepsilon m}{1+\varepsilon} - n \ge \varepsilon m/2.$$

Let  $\gamma' = \lfloor ip + \varepsilon np \rfloor$  and  $\gamma = \lceil ip + 2\varepsilon np \rceil$ . Let us choose  $y_1, \ldots, y_{\gamma} \in I$ , one by one, independently and uniformly distributed. Of course, some of these might coincide. Let  $Y := \{y_1, \ldots, y_{\gamma}\}$ , ignoring multiple occurrences of the same vertex. The probability that  $|Y| \le \gamma'$  is at most

$$\begin{pmatrix} i \\ \gamma' \end{pmatrix} \left( \frac{\gamma'}{i} \right)^{\gamma} \leq \left( \frac{\mathrm{e}i}{ip + \varepsilon np} \right)^{ip + \varepsilon np} \left( \frac{ip + \varepsilon np}{i} \right)^{ip + 2\varepsilon np}$$
$$= \mathrm{e}^{ip + \varepsilon np} \left( p + \frac{\varepsilon np}{i} \right)^{\varepsilon np}.$$

The last expression, as a function of a real-valued argument  $i \ge 0$ , is first decreasing and then increasing in *i* so it is maximised if either i = n or *i* achieves the lower bound (14). In either case, the result can be bounded by  $o(n^{-1})$ . Thus, the set *Y* has at least  $\gamma'$  elements with probability  $1 - o(n^{-1})$ .

Let the random variable U count the number of  $x \in S$  which belong to none of K - l(y),  $y \in Y$ .

We consider the martingale  $(U_0, \ldots, U_{\gamma})$ , where  $U_j$  is the expected value of U after having exposed the first j vertices  $y_1, \ldots, y_j$ . Clearly, each new vertex changes U by at most k.

It is easy to estimate  $U_0$ , the expectation of U:

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$$E(U) \ge |S| \frac{\binom{l-i}{\gamma}}{\binom{i}{\gamma}} \ge s \left(\frac{i-t-\gamma+1}{i-\gamma+1}\right)^{\gamma} \ge s e^{-(1+\varepsilon)\gamma t/i} \ge s(np)^{-\frac{1}{2}+\frac{\delta}{4}}.$$

(Note that t = o(i) by the definition of t and  $\gamma = o(i)$  by (14).)

By applying the Hoeffding–Azuma inequality [2, 11] (see e.g. Alon and Spencer [1, Theorem 7.2.1]) we obtain

$$Pr\{U=0\} \le Pr\{|U-E(U)| \ge E(U)\} \le \exp\left(-\frac{(E(U))^2}{2k^2\gamma}\right)$$
$$= \exp(-\Omega((np)^{\frac{\delta}{2}}/\ln^2(np))) = o(n^{-1}).$$

Hence, the event that  $|Y| < \gamma'$  for some *i* or U = 0 has probability o(1). Of course, when we select a random *a*-subset of *Y*, we obtain a uniformly distributed *a*-subset of *I*. Note that  $|\Gamma(i + 1) \cap [i]| \le \gamma'$  by (11). We can find a distribution for  $a \in [0, i]$  such that when we first choose *a*, then *Y* as above, then a random *a*-subset of *Y*, we obtain precisely the distribution of  $\Gamma(i + 1) \cap [i]$ , conditioned on (11).

Hence, almost surely for any *i*, condition (13) holds; that is, we can always choose an appropriate label.  $\Box$ 

#### 4. General graphs

Let us prove upper bounds that apply to arbitrary graphs. The obvious greedy algorithm gives the following (cf. Wood [15, Theorem 4]).

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Lemma 6. For any graph G we have

$$\mathcal{D}(G) \le 2\Delta(G)e(G) + v(G).$$
  
$$\mathcal{S}(G) \le \Delta(G)e(G) + v(G).$$

**Proof.** Let us bound S(G), for example. We choose vertex labels one by one. When we consider a vertex  $i \in V(G)$ , we are forbidden to choose a previously used label as well as any number of the form l(u) + l(v) - l(w) where  $\{u, v\}, \{w, i\} \in E(G)$  and the labels of u, v and w have already been chosen. This forbids at most v(G) - 1 + d(i)e(G) elements so we can always proceed.  $\Box$ 

Remarkably, the trivial Lemma 6 is not far from the truth: if applied to  $G \in \mathcal{G}(n, p)$ , with  $\frac{np}{\ln n} \to \infty$ , it gives a bound within the multiplicative factor of  $O(\ln n)$  from the actual value. It is an interesting open problem to determine the maximum value of  $\mathcal{D}(G)$  and  $\mathcal{S}(G)$  over all graphs of order *n* and maximum degree at most *d*.

For the functions  $\mathcal{D}(n, m)$  and  $\mathcal{S}(n, m)$  we obtain the following upper bounds.

**Theorem 7.** Let  $n \to \infty$  and  $m \leq {n \choose 2}$ . Then

$$\mathcal{D}(n,m) \le n + (2^{4/3} + o(1))m^{4/3},$$
  
$$\mathcal{S}(n,m) \le n + (2^{2/3} + o(1))m^{4/3}.$$

**Proof.** Let us deal with  $\mathcal{D}(n, m)$  here. Let *n* be large and *G* be an arbitrary graph of order *n* and size *m*. It is easy to see that  $\mathcal{D}(n, m) = n$  if m = O(1), so assume that  $m \to \infty$ .

Order the vertices of *G* by their degrees:  $d(x_1) \ge \cdots \ge d(x_n)$ . Let  $k = \lceil (2m)^{2/3} \rceil$ . Label vertices  $x_1, \ldots, x_k$  by a Sidon *k*-subset of  $[s], s = \lfloor (1 + \varepsilon)k^2 \rfloor$ . We try to label the remaining vertices one by one using labels from [n + s]. When choosing a label for  $x_i$ , the forbidden values are the already assigned labels, i - 1 of them, as well as the numbers  $l(x_u) \pm (l(x_v) - l(x_w))$ , where  $u, v, w \in [i - 1]$  with  $\{x_i, x_u\}, \{x_v, x_w\} \in E(G)$ , at most  $2md(x_i)$  numbers. But  $d(x_j) \ge d(x_i)$  for any j < i, hence  $d(x_i) \le \frac{2m}{i} \le \frac{2m}{k}$  and the total number of forbidden labels is at most

$$n-1+\frac{4m^2}{k} < n+s;$$

that is, we can always find a suitable label.  $\Box$ 

Needless to say, we have a trivial upper bound, namely  $(1 + o(1))n^2$ .

Good lower bounds on  $\mathcal{D}(n, m)$  and  $\mathcal{S}(n, m)$  are provided by random graphs plus isolated vertices. Our aim is to choose  $v \le n$  such that, if we define p by  $p\binom{v}{2} = (1-\varepsilon)m$ , the bound of Theorem 4 for  $G \in \mathcal{G}(v, p)$  is as large as possible. In order not to clutter this paper with details we compute only the order of magnitude, not bothering about multiplicative constants.

If  $m = \Omega(n^{3/2}\sqrt{\ln n})$ , then we take v = n. Almost surely  $\mathcal{D}(G)$ ,  $\mathcal{S}(G) = \Omega(n^2)$ . Otherwise, take  $v = \Theta(m^{2/3}(\ln m)^{-1/3}) < n$ . Now, the lower bound is

$$\mathcal{D}(n,m), \mathcal{S}(n,m) = \Omega(m^{4/3}(\ln m)^{-2/3}), \quad \text{for } m = o(n^{3/2}\sqrt{\ln n}).$$

Also, note the trivial lower bound  $\mathcal{D}(n, m)$ ,  $\mathcal{S}(n, m) \ge n$ .

A little more careful analysis shows that there is an absolute constant *C* such that our lower and upper bounds on  $\mathcal{D}(n, m)$  and  $\mathcal{S}(n, m)$  are within factor  $C(\ln n)^{2/3}$  for any *m*, *n*. This poses an intriguing problem of closing this gap.

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